TCP/IP NETWORK ROUTE PERFORMANCE ANALYSIS

Including a detailed look at how AES CLEVER *e*Route 3.0 analyzes network routes and identifies the root-causes of performance problems

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TCP/IP NETWORK ROUTE PERFORMANCE ANALYSIS

INTRODUCTION

Burgeoning mission-critical applications are heavily dependent on Internet technologies. Users demand more and more functionality, with near zero downtime. These ever-increasing demands make the management of enterprise networks extremely challenging. When there is a response time issue, the user needs to know where to look for information and then exactly what to look for in order to resolve the problem. Most networks have a limited set of shared resources and there are always potential performance bottlenecks elsewhere in the network. It is counter-productive to simply eliminate a performance bottleneck, only to have it move to some other resource.

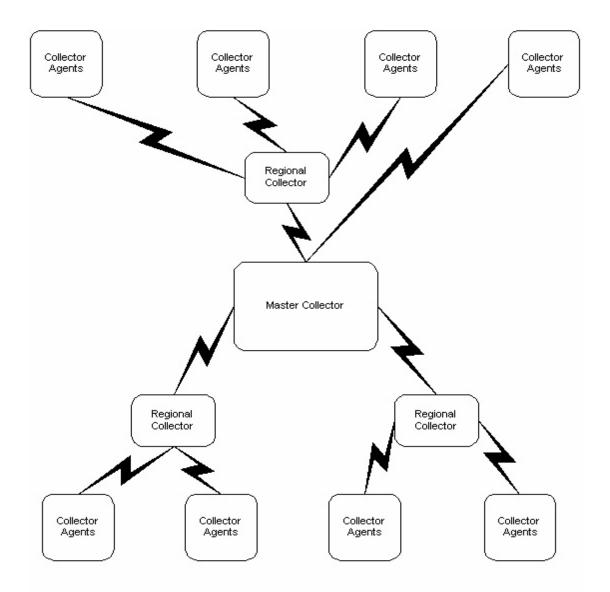
Internet performance can cover a wide variety of components, each of which can affect overall network performance. Applications, Web servers, Intranets, Local Area Networks (LANs), protocols, operating systems, and Wide Area Networks (WANs) can have an impact on overall Internet performance. This paper discusses only the network-related performance concerns, such as TCP/IP routing analysis.

The Hyper Text Transfer Protocol (HTTP) is an application-layer protocol sitting on top of the TCP/IP to establish the communication path between a Web surfer and Web server. HTTP maps the server name to an IP address, establishes a communication path with the Web server, and then sends the request using the Uniform Resource Locator (URL). It also receives the response from the Web server (e.g., Hyper Text Markup Language or imaging document). When the communication is complete, HTTP closes the TCP/IP connection. One needs to understand the potential performance impact of the HTTP and TCP/IP. For example, the protocol latency delay may consist of path delay and processing delay. Path delay refers to the time required to establish a path between the Web surfer and the Web server. The process delay refers to the request processing time across the established path.

Analyses of TCP/IP routes help network administrators and performance analysts identify the root-causes of performance problems. For example, a problem that may appear to be caused by poor transaction response time may actually be caused by network route congestion, gateway failure, packet loss, or an unreachable destination, making performance data collection just as important as data analysis. Without a good set of data with statistical significance, the results of analyses may be misinterpreted and problems may not be resolved correctly. This could result in unnecessary tuning or capacity upgrades.

PERFORMANCE DATA COLLECTION ARCHITECTURE

Transaction failure, transaction performance, and network resource usage data can be useful for analyzing TCP/IP transactions. Analyses are performed at the TCP/IP session level. One possible implementation of data collection architecture is a three-layer hierarchy that consists of a centralized Master Collector & Processor, regional Collectors, and Remote Collector Agents. The diagram below shows the structure.



Each collector agent gathers TCP/IP transaction data using the TraceRoute command on a set of network nodes assigned under it. The Command transmits packets to each network node that constitutes a path or route between the data collector and the network node. The response time information on that path or route, along with the number of segments (also referred to as hops, the path between two network nodes), and the host node detected along the path will be collected and saved. The remote data collector agents transmit saved data at predetermined time intervals to their regional collectors. The regional collectors then transmit the gathered performance data to the Master Collector. The Master Collector can process the raw data and produce either scheduled or ad hoc reports.

3-Layer Data Collection Architecture

Three-layer architecture is somewhat complex for most implementations, however, and the collection model can be simplified and reduced to two layers in even the largest networks, whereby the Master Collector can directly control the remote data collection agents. A typical application of such architecture might be to install the Master Collector on a Windows system

and remote data collection agents on Windows/LINUX systems, with the master collector also providing the analysis function. Multiple Master Collectors can be implemented in such a scenario, if necessary, making the two-tier structure ideal. CLEVER *e*Route, developed by AES, is an extremely good implementation of the two-layer architecture just described.

Three categories of data provide detailed insight, enabling network analysts to identify the cause(s) of the problem or the resource affected by a performance bottleneck:

- Transaction Failure data provides the times and types of failures, such as packet loss, looping, or connection failure, for each route or segment.
- Transaction performance, analyzed per segment or route, tracks the response time distribution over a selected time period, including peaks and valleys of transaction volumes.
- Network resource usage gives an indication of productive and non-productive use of network resources with detected peaks and valleys.

Network analysts can quickly isolate failed routes and segments by viewing these categories of performance data along with the overall network topology. Transaction volume statistics, such as patterns of shared resource usage dominance, can also be useful for better managing future network capacity growth.

One needs to be aware that TCP/IP route/path analysis only covers one aspect of overall TCP/IP performance analysis. It can be used in conjunction with other types of tools such as Web Performance analysis of selected URL pages or Web page usage trending analysis, to address specific areas of concern.

PERFORMANCE DATA COLLECTION

A solid data collection system design is vital to the success of subsequent performance analysis. There are four basic procedures for data collection and analysis:

- Configuration the master data collector can configure all remote collector agents by issuing a Configure Command. Any remote data collector agent can be re-assigned to a different master data collector, if present.
- Collection the master data collector can issue either real-time or periodic scheduled collections to data collection agents, each of which collects data locally on those network nodes where a collection module is initiated.
- Analysis as the collection module is invoked, collected data are stored in the application's subdirectory for identification purposes. The collection module also imports, analyzes, and generates reports from the collected data.
- Identification each physical network device can be identified via the device's Simple Network Management Protocol (SNMP) Management Information Base (MIB) information such that data can be correlated to the physical device.

Depending upon where and how the data collection mechanism is set up, users can select certain data collection parameters. The TraceRoute Command can pick up the Maximum Time to Live (TTL) or Number of Maximum Network Hops value for the analysis. The Timeout value in seconds specifies how long to wait for a response. The Interval value in minutes defines the time between collecting two consecutive samples. The Sample value determines the total

TCP/IP Network Route Performance Analysis

numbers of samples that need to be collected. Users can select a particular time interval on each day to collect TCP/IP route data or can run the data collection continuously. There is no ideal data-sampling interval or number of samples since each networking environment has different characteristics. Data could be collected for example several times a minute, or every five minutes, and summarized every 15 minutes.

PERFORMANCE DATA REPORTING

(a) For route or Segment Analysis Reporting

Since a route failure may caused by several factors, analysis reports should cover, as a minimum, the following information, from which defective routes, segments and devices can be further identified:

- Session identifier assigned to the session during the collection period
- Route identifier assigned to the route during the analysis
- Target/Initiating Host as the starting point of the route including its IP address
- Receiving/ending Host of the router including its IP address
- Number of times the route was analyzed and reported
- Date on which the route failure occurred
- Time at which the route failure occurred
- Segment identifier assigned to the segment during the analysis & reporting period
- Response time in milliseconds the time to takes to perform a Trace Route Command from the initiating host to the receiving host
- Data Collection time (time at which the route/segment information was collected)
- Data Collection date (date on which the route/segment information was collected)
- Peak/Valley times, including the minimum, average and maximum response times based on the total number of samplings collected for the route. The average response time is calculated by averaging the total response time of all hops preceding the current hop and subtracting that from the current hop's response time.
- Number of successful segments in the route under analysis
- Number of failed segments in the route under analysis
- Percentage of segment availability (the ratio between successful and failed segments)
- Status (route/segment operating status for availability reporting)
- Failure device type, based on the highest failing, looping, or packet losing routes

(b) SNMP MIB Data Analysis Reporting

It is worthwhile to note that SNMP was designed to accommodate different network platforms, various protocols, and proprietary operating systems. SNMP exchanges network information through messages known as Protocol Data Units (PDUs). There are five useful PDUs for network monitoring: two deal with setting terminal data, two deal with reading terminal data, and one (called trap) is used for monitoring network events such as start-up and shut-down. There are two major types of MIBs: public and private. Public MIBs contain information that most network vendors support and private MIBs contain network vendor-specific information.

Reports that utilize public MIB information should contain at a minimum the following information:

Border Gateway Protocol (BGP)	Used for exchanging routing infor each with its own route in a netwo should contain the BGP version, I number, and BGP identifier of the also contain a table of its peer BG	ork of autonomous systems. It ocal autonomous system local system. Each BGP should
	 Local IP Address and Port Remote IP Address, Port, an Messages-In Messages-Out Connection Retry Peer BGP's Holding Time Re Keep-Alive Time Time-To-Live (TTL) Values 	
CSS Bridge	Contains a list of information for e contain the following information:	each port of a bridge. It should
	 kept Interface Index for the interfa Port Circuit which is used to a by another port on the same Delay-Exceeded Discards (the this port due to excessive transport on the same transport on the same transport due to excessive tran	distinguish the same value used bridge he number of frames discarded by nsit delay through the bridge) e number of frames discarded by
CSS Ethernet	Contains statistics for interface to should contain such information a	
	Index	A unique identifier
	Alignment Errors	The count on frames received without an integral number of octets in length
	Frame Check Sequence (FCS)	Refers to the extra characters added to a frame for detecting and correcting errors
	Errors	The count of frames due to FCS check

Single Collision Frames	A count of successfully transmitted frames due to a single collision
Multiple Collision Frames	A count of successfully transmitted frames that have more than one collision
Deferred Transmissions	A count of frames that experienced delay not due to collision
Excessive Collisions	A count of frames whose transmission failure was due to excessive collisions
Carrier Sense Errors	The number of times that the carrier sense condition was lost or never asserted when attempting to transmit a frame
Frame Too Long	A count of frames that exceed the maximum permitted frame size
Internal MAC Receive/Transmit Errors	A count of frames failed due to either MAC sub-layer receiver/transmitter error
CSS RMON Contains remote network monitor contain such information as the features of the second se	ring statistics information. It should ollowing:
 are dropped due to lack of re Octets (total number of octet) Packets (total number of packets (total number of packets) Undersize Packets (<64 octet) Oversize Packets (>1518 oc) Fragments (total number of packets than 64 octets with a backets) Collisions (total number of examples) 	er of events that probe packets esources) is of data) ckets including bad, broadcast and ets) tets) packets received whose length is

Status

Internet Control Message Protocol (ICMP)	and received by the loca received in error, or not s should contain the follow	p track of the message types generated I ICMP entity including received, sent, sent due to error. The basic information ving:
	 Packet-In Packet-Out Errors Timing (time stamps ICMP Address Mask 	s, replies) < Messages (requested & received)
Interface	Contains generic information	ation on the interface layers, such as:
	 Interface Description Type Maximum Transmission Rate (Media Specific Addr Administrative and O Discards (IN & OUT Out Queue Length 	sion Unit (MTU) size (speed, number of INs & OUTs) ess Operational Status
IP MIB	(e.g., if the device is acti	the IP subsystem of a managed node ng as a router or a host, packets or arded, time-out values, etc.). It should ation as:
		ed, discarded) warded, discarded) ress, unknown protocol, no routes) success, failed, time-out)
IP Route	Contains the following in	formation:
	Destination address	Destination address, on this route and various routings, for primary routing, four alternative routes, Next Hop, etc.
	Protocol	Learning various route protocols on possible four values: none, local, network management or obtained via ICMP
	Age	Time in seconds since last route update
	Mask	For comparison of destination address
	Reference	Specific to the particular routing protocol responsible for this route

Open ShortestUsed with large autonomous networks in preference to the RoutingPath First (OSPF)Information Protocol (RIP). It should contain:

Router information	A 32-bit integer for the router in the autonomous system, and whether it is an area border router
OSPF Status	Version, administrative status
Link State Advertisements (LSA)	External LSA, Number of Originated & Received new LSAs, Exit Overflow Interval, etc.

Within OSPF, there are six sub-groups; each providing further information on OSPF performance:

OSPF Area	Contains a table with information on the configuration parameters and cumulative statistics of a router's attached areas. Information such as the following should be included:
	 Authentication Types Imported Autonomous Systems External Link State Reachable Board Routers Area Status
OSPF Link State Data Base	Contains a table with information about the link state announcement from attached neighboring areas. Information such as the following should be included:
	 Area IDs Type Originating Router ID Age Checksum LSA Content
OSPF Host	Contains a table with information on directly attached hosts to the router. Information such as the following should be included:
	 Host IP address Type of Service Status Area ID to which the host is

OSPF Interface	Contains a table with information describing interfaces on this OSPF. Information pertinent to interface analysis includes the following: - Interface IP address - Area ID - Type - Status - Priority - polling interval - authentication type - key
OSPF neighbor	Contains a table with information on all neighbors in the locality of this router. Information important to the analysis includes the following: - neighbor ID - router IP address - priority - operating status
OSPF Area Advertiseme nt Table	Replaces the OSPF Summary table containing the following Area Advertisement information: - ID - Type - IP Addresses - Subnet Mask - Effect (advertise matching or not matching)
printer. A critical alert is o	non-critical alerts currently active in one that stops the printer from printing ther printing. Non-critical alerts are

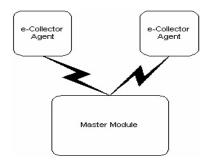
Printer Alert Contains a list of critical and non-critical alerts currently active in the printer. A critical alert is one that stops the printer from printing immediately and prevents further printing. Non-critical alerts are those that do not stop printing, but may become critical at a future time. This table should contain information on severity, component and its detailed location, alert code, and a description of each currently active critical alert.

ТСР	The Transmission Control Protocol MIB information identifies the
	retransmission algorithm, maximum and minimum retransmission
	time-outs, and number of active/passive opens, resets, connects,
	etc. Information important to the analysis includes the following:

- Segment-In/Out
- Retransmission
- Numbers of Attempts
- Opens
- Resets
- Maximum Connections
- Time-Out Value (Min and Max)
- **UDP** User Datagram Protocol (UDP) provides four counters and a table for datagram delivery, destined for unknown ports, discarded due to format errors, and send from UPD group. Information important to the analysis includes the number of total datagrams delivered/sent, and errors types.

APPLICATIONS & EXAMPLES

A two-layer architecture of the data collection system, CLEVER *e*Route, is used to demonstrate the application of TCP/IP route analysis. A Master Module can be installed on a workstation where the data collection commands and reports are generated. A number of data collection agents, or eCollectors, are installed on remote Linux or Windows machines. The eCollectors, in addition to performing route information collection, are also capable of auto-discovering local IP addresses, executing run-time scripts, launching the collection scheduler, and maintaining data logs.



The 2-LayerArchitecture Employed by CLEVER eRoute

CLEVER eRoute obtains its data from the eRoute Host Collector (if present), the eCollector (local to the Master), and the remote eCollector Agents. Each eCollector Agent performs a TraceRoute command from the remote starting point to a set of network nodes (any IP-addressable device). The TraceRoute command transmits packets to each node along the path

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that constitutes the route and outputs a millisecond response time for each packet. Each response time associated with each network node is collected and saved in an output file that provides raw data for CLEVER *e*Route to analyze. The output file provides information on route performance, number of segments (the path between two network nodes, also called hops), and each host detected along the route. From this information CLEVER *e*Route analyzes the data, systematically organizing and presenting it to the user in an accessible and easily understandable format. Among its unique features, CLEVER eRoute has the ability to collect performance statistics *from* anywhere in the network *to* anywhere in the network. Its eCollectors can run outbound and inbound simultaneously, facilitating the analysis of return routes.

Analysis Reports

The following analysis reports are available for review:

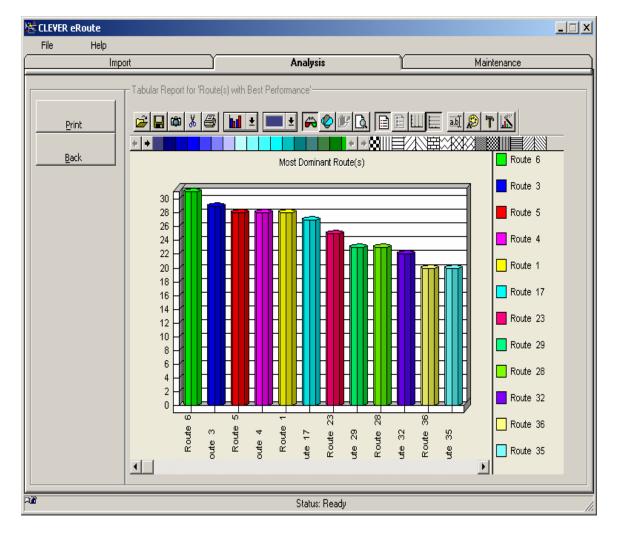
- Most Failing Route
- Most Failing Segment
- Failures by Time of Day
- Packet-Loss by Time of Day
- Most Packet-Losing Route
- Most Packet-Losing Segment
- Looping by Time of Day
- Most Looping Route
- Most Looping Segment
- Route with Best Performance
- Segment with Best Performance
- Route with Worst Performance
- Segment with Worst Performance
- Performance by Time of Day
- Route Performance Distribution
- Segment Performance Distribution
- Most Dominant Route
- Most Shared Host(s)
- Most Shared Segment(s)

TCP/IP Network Route Performance Analysis

File Help	ort	Analy	zis.	Maintenance	2
				indiritoridirioo	
	Route Analysis				
Select Session(s)	Starting Points	Ending Points	Analysis Report	Graphical Report	Tabular Report
201000 2022/01/07	- Session#5		Most Failing Route(s)		
Performance Profile		208.254.23.156	Most Failing Segment(s)		
	ga2 ga2	66.150.174.65 216.239.35.100	Failure(s) by Time of Day		
Route Analysis		204.202.131.230	Most Packet-Losing Route(s)		
an and an a	- Session#4		Most Packet-Losing Segment(s)		
Segment Analysis	INTSMVS	64.139.36.116	Packet Losses by Time of Day		
	INTSMVS INTSMVS	64.139.36.117 208.55.14.112	Most Looping Routes		
Route Statement	INTSMVS	66.218.71.95	Most Looping Segments		
egment Statement	- Session#1		Looping(s) by Time of Day		
egment Deatement	qa3	www.l1p.net	Route(s) with Best Performance		
Defective Routes	qa3	www.dwyt.com	Segment(s) with Best Performance		
	ga3 ga3	207.68.171.233	Route(s) with Worst Performance		
efective Segments			Segment(s) with Worst Performance		
			Route Performance Distribution		
			Segment Performance Distribution		
			Performance by Time of Day		
			Most Dominant Route(s)		
			Most Shared Host(s)		
	Reporting Top 10	Reports of Selected Item	Most Shared Segment(s)		

CLEVER eRoute Network Route Analysis

The Route Analysis menu shows two sections of information. The first one provides the selection options for display (e.g., performance profile of the selected session, route segments, segment statement, defective routes and defective segments). The starting and ending points for the selected sessions are listed for selection. The second one, the right side table, lists all the available reports with options for graphical or tabular formats. An example of a graphic format report is shown below:



An Example of CLEVER eRoute Graphic Reports

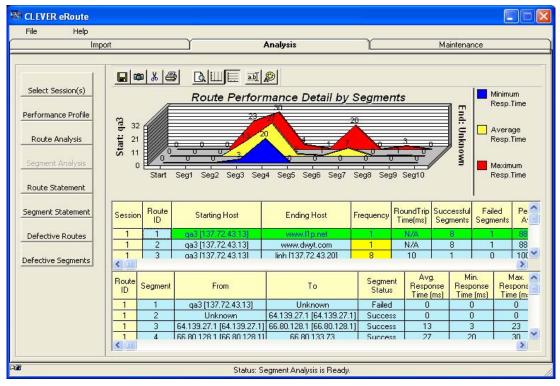
The appropriate tabs can be selected to view the best and worst routes and segments. The report contains the following information:

- Session the identifier assigned to the session during the data collection period
- Route ID the identifier assigned to the route during the data collection period
- Segment ID the identifier assigned to this segment during data collection period
- Response time the time it takes to perform a TraceRoute Command from the starting point to the ending point, in milliseconds
- Frequency the number of times the route was reported
- Collect Time the time at which the route information was collected
- Collect Date the date on which the route information was collected
- Starting Host starting point for the route
- Ending Host ending point for the route

The report can show the minimum, maximum, and average response time to illustrate Performance by Time of Day:

- Minimum RT (ms) the minimum response time for the segment based upon the total number of data scans performed for the route
- Maximum RT (ms) the maximum response time for the segment based upon on the number of data scans performed for the route
- Averaged RT (ms) the averaged response time for the segment, which is calculated by averaging the total response times of all network hops preceding the current one and subtracting that from the current hop's response time

A Segment Analysis report for a given segment can be displayed to provide detailed performance data. Again, the display can be in either graphical or tabular format.



CLEVER eRoute Route Performance Detail by Segment

The sample graphic display report shows two sections. The upper section is a summary of the route information for the selected session with minimum, averaged and maximum response times. The lower section shows the following detailed segment performance data:

- Session ID
- Route ID
- Staring Host
- Ending Host
- Frequency
- Round Trip Time (ms) the amount of time it took for the packet to reach its Ending Host and return to its Starting Host (if the route fails, the display will indicate N/A)
- Successful Segments the number of segments that succeeded in the route
- Failed Segments the number of segments that failed due to unavailability of the next IP address in the amount of time specified for the data collection

- Percent Availability the ratio between successful and failed segments
- Status displays complete or partial, depending upon the successful completion or failure of the route

In the MIB data analysis, there are several very useful reports for route performance analysis.

CSS-Ethernet Report

Some of the data fields in the CSS-Ethernet report (see example below) can provide insight into Ethernet interface performance. The Single Collision Frames column provides the count of successfully transmitted frames on a particular interface for which transmission is inhibited by more than one collision. The Late Collision field shows the number of times that a collision is detected on a particular interface later than 512 bit-times into the transmission of a packet. The Excessive Collisions field points towards a particular interface failing due to excessive collisions. The Frames Too Long data field shows the count of frames received on a particular interface that exceed the maximum permitted frame size.

ndex	Alignment Errs	FCS Errs	Single Collision Frames	Multiple Collision Frames	SQE Test Errs	Deferred Transmissions
9	1089	1067	1097	1113	1057	1129
2	1103	1145	1098	1040	1137	1099
3	1017	1111	1087	1169	1071	1184
4	1084	1145	1136	1172	1172	1075
5	1100	1191	1073	1096	1067	1070
6	1079	1041	1069	1144	1078	1100
7	1118	1007	1080	1141	1141	110-
8	1150	1151	1056	1036	1116	117
9	1118	1160	1068	1074	1022	114
10	1160	1108	1049	1093	1096	116

ndex	Late Collisions	Excessive Collisions	InternalMac Transmit Errs	Carrier Sense Errs	Frames TooLong	InternalMac Receive Errs
ä	1015	1180	1159	1144	1069	113:
2	1155	1108	1045	1090	1111	113
3	1169	1090	1112	1060	1096	116
4	1183	1162	1182	1111	1118	112
5	1087	1072	1167	1088	1103	101
6	1051	1033	1090	1142	1131	106
7	1095	1086	1139	1172	1113	103
8	1171	1133	1162	1116	1106	118
9	1084	1072	1161	1066	1095	104
10	1050	1106	1108	1112	1154	111
0						

CLEVER eRoute CSS Ethernet Reporting

CSS – RMON Report

The CSS RMON report provides remote network monitoring statistics, as seen in these examples:

ndex	Data Source	Drop Events	Octets	Packets	Broadcast Packets	Multicast Packets	FCS-Align Errors
4	1.2.3	1079	1014	1072	1042	1072	1042
2	1.2.3	1028	1008	1030	1049	1012	109
3	1.2.3	1087	1070	1073	1070	1002	105
4	1.2.3	1079	1071	1036	1062	1067	109
5	1.2.3	1019	1085	1020	1043	1079	104
6	1.2.3	1020	1053	1090	1041	1068	106
7	1.2.3	1025	1084	1029	1003	1047	101
8	1.2.3	1026	1022	1008	1004	1014	101
9	1.2.3	1081	1058	1097	1025	1001	100
10	1.2.3	1070	1079	1079	1088	1088	108
							8

Index	Undersize Packets	Oversize Packets	Fragments	Jabbers	Collisions	Packets64octets
4	1060	1029	1055	1009	1052	1075
2	1012	1040	1076	1047	1030	1062
3	1020	1070	1070	1056	1069	1013
4	1087	1020	1096	1098	1047	1042
5	1042	1042	1040	1010	1053	1039
6	1031	1072	1030	1067	1091	1080
7	1007	1017	1081	1037	1067	105
8	1090	1079	1073	1034	1068	108
9	1012	1057	1045	1016	1086	1017
10	1070	1053	1015	1073	1085	1009
<			(mil)			>

ndex	Packets64octets	Pkts64to127	Pkts128to255	Pkts256to511	Pkts512to1023	Pkts1024to1518	Owner	Status
1	1167	1101	1103	1152	1080	1263	abc	1
2	1148	1210	1202	1083	1100	1188	abc	1
3	1114	1194	1129	1238	1188	1137	abc	1
4	1122	1151	1176	1215	1165	1130	abc	1
5	1141	1093	1089	1132	1165	1111	abc	1
6	1177	1121	1150	1185	1133	1132	abc	1
7	1109	1247	1118	1201	1176	1108	abc	1
8	1150	1088	1118	1152	1120	1110	abc	1
9	1129	1264	1160	1115	1196	1084	abc	1
10	1142	1204	1236	1108	1158	1145	abc	1

CLEVER eRoute CSS RMON Reporting

The Number of Octets field shows the total number of octets received on the network, and can be used to estimate Ethernet utilization. The Packet information (broadcast, multicast, undersized, oversized), and the brackets of packet sizes (up to 64 bytes, 65-127 bytes, 128-255 bytes, 256-511 bytes, 512-1023 bytes, 1024-1518 bytes) can provide packet volume and distribution information.

ICMP MIB Report

The Internet Control Message Protocol (ICMP) objects are the input and output error and control message statistics for the IP layer. The report groups packets by IN and OUT categories, and provides a graphical report format.

	<u>In</u>	Out	
Packets:	5,481	4,886	
Errors:	1,521	1,842	Packets
Dest. Unreachable:	1,615	1,551	Errors
Time Exceeded:	1,469	1,546	DestUnrch
arameter Problems:	1,690	1,528	
Source Quench:	1,417	1,602	TimeExcd
Redirects:	1,586	1,827	ParmProb
Echos:	1,548	1,711	SrcQuenct
Echo Replies:	1,785	1,575	Redirect
Timestamps:	1,573	1,736	
Timestamp Replies:	1,722	1,633	Echos and a second
Address Mask:	1,549	1,748	EchoRep -
ress Mask Replies:	1,609	1,624	Timestmp
Print	0K		TimestmpF
uto-Refresh			AddrMask
Address Mask: Iress Mask Replies:	1,549 1,609	[1.748 [1.624	Timestmp Timestmpf

CLEVER eRoute ICMP MIB Reporting

The Packets data field provides the total number of ICMP packets received and sent. The Errors field shows the number of receiving/sending ICMP errors (e.g., bad ICMP checksum, bad length, inability of IP to route, etc) used for determining causes of errors. The number of time-outs and time-stamps gives an indication as to the number of ICMP request and reply messages received and sent for coordinating the transmission sequence.

ndex	In Non UCast	In Discards	In Errors	In Unk. Prot	Out Octets	Out UCast	Out Non UCast	Out Discards	Out Errors	Out Q Len	IF Specific OID
1	1,684	1,628	1,433	1,507	1,826	1,751	1,588	1,562	1,586	907	1.2.3
2	1,620	1,568	1,736	1,642	1,487	1,554	1,472	1,690	1,633	1,041	1.2.
3	1,691	1,670	1,888	1,610	1,831	1,638	1,682	1,767	1,713	936	1.2.
4	1,575	1,615	1,855	1,601	1,595	1,646	1,610	1,577	1,560	1,291	1.2.
5	1,625	1,709	1,537	1,458	1,576	1,516	1,819	1,564	1,620	994	1.2.
6	1,782	1,589	1,665	1,472	1,515	1,492	1,613	1,563	1,606	1,056	1.2.
7	1,703	1,557	1,724	1,513	1,558	1,627	1,573	1,534	1,507	720	1.2.
8	1,789	1,761	1,672	1,665	1,422	1,675	1,745	1,513	1,539	1,245	1.2.
9	1,578	1,597	1,675	1,604	1,579	1,585	1,511	1,904	1,739	1,252	1.2.
10	1,459	1,597	1,739	1,537	1,581	1,360	1,509	1,538	1,683	981	1.2.
6											

CLEVER eRoute Interface Reporting

Interface Report

The Interface report contains information regarding the entity's interfaces. Each interface is treated based on the pertinent sub-network.

ndex	Description	Туре	MTU	Speed	Phys. Address	A-Status	0-Status	LastChg.	In Octets	In UCast
1	My_ifDescr_string	Other	1	1,029	979899	Up	Up	6577100	1,764	1,76
2	My_ifDescr_string	Other	1	1,257	979899	Up	Up	6577100	1,546	1,41
3	My_ifDescr_string	Other	1	906	979899	Up	Up	6577100	1,939	1,63
4	My_ifDescr_string	Other	1	986	979899	Up	Up	6577100	1,740	1,74
5	My_ifDescr_string	Other	1	976	979899	Up	Up	6577100	1,745	1,67
6	My_ifDescr_string	Other	1	822	979899	Up	Up	6577100	1,698	1,77
	My_ifDescr_string	Other	1	1,072	979899	Up	Up	6577100	1,695	1,49
	My_ifDescr_string	Other	1	1,475	979899	Up	Up	6577100	1,638	1,73
9	My_ifDescr_string	Other	1	594	979899	Up	Up	6577100	1,556	1,57
10	My_ifDescr_string	Other	1	1,042	979899	Up	Up	6577100	1,819	1,75
(100)		4								

ndex	In Non UCast	In Discards	In Errors	In Unk. Prot	Out Octets	Out UCast	Out Non UCast	Out Discards	Out Errors	Out Q Len	IF Specific OID
1	1,684	1,628	1,433	1,507	1,826	1,751	1,588	1,562	1,586	907	1.2.
2	1,620	1,568	1,736	1,642	1,487	1,554	1,472	1,690	1,633	1,041	1.2.
3	1,691	1,670	1,888	1,610	1,831	1,638	1,682	1,767	1,713	936	1.2.
4	1,575	1,615	1,855	1,601	1,595	1,646	1,610	1,577	1,560	1,291	1.2.
5	1,625	1,709	1,537	1,458	1,576	1,516	1,819	1,564	1,620	994	1.2.
6	1,782	1,589	1,665	1,472	1,515	1,492	1,613	1,563	1,606	1,056	1.2.
7	1,703	1,557	1,724	1,513	1,558	1,627	1,573	1,534	1,507	720	1.2.
8	1,789	1,761	1,672	1,665	1,422	1,675	1,745	1,513	1,539	1,245	1.2.
9	1,578	1,597	1,675	1,604	1,579	1,585	1,511	1,904	1,739	1,252	1.2.
10	1,459	1,597	1,739	1,537	1,581	1,360	1,509	1,538	1,683	981	1.2.

CLEVER eRoute Interface Reporting

The Type of interface is based on the protocol immediately below the network layer in the protocol stack. The Maximum Transmission Unit (MTU) limits the size of the largest datagram that can be sent/received on the interface. The Speed is an estimate of the interface's current bandwidth in bits per second. This table also provides a set of causes for IN and OUT packets of the interface (e.g., Discards, Errors, Unsupported, Non-Unicast, etc.). The OUT Q Len field indicates the length of the output packet queue, which is very helpful in determining if the output queue size is properly chosen. The total numbers of IN and OUT packets in octets are a good indication of the performance of the interface.

IP MIB Report

The data fields of the IP MIB report, shown below, provide several useful statistics.

Forwarding:	Gateway		
Default TTL:	1	Receives	
In Receives:	1	Delivers	
Delivers In:	1,620	in in in its second	
Requests Out:	1,524	Out	
Datagrams Forwarded:	1,316	Dgrms	
Header Errors In:	4,757	Hdr Errs	
Address Errors In:	1,550	in .	·
Unknown Protocol In:	1,559	Addr Err	
Discards In:	1,532	Unk. Prot	
Discards Out:	4,730		
No Routes Out :	1,513	In	
Reassembly Timeout:	1,569	Discards Out	
Reassembly Required:	1	No Route	~
Reassembly OK:	1.511		
Reassembly Failed:	1,421	0 681 1362 2043 2724 3405	4086 4767
Fragment OK:	1,416		[2000 <u>000000000000000000000000000000000</u>
Fragment Failed:	1,514	Print OK	Route Ta
Fragments Created:	1 010	Auto-Refresh	

CLEVER eRoute IP MIB Reporting

The Time-To-Live (TTL) field in the IP header of datagram shows the TTL value, ranging from 1 to 255. The Forwarding field indicates whether this entity is acting as an IP gateway. The value 1 indicates to forward the received packet and a 2 means do not forward. The Reassembly information, such as the number of requests, time-outs, failed and successful counts, would provide the effectiveness of the datagram reassembly process. The Fragment field gives an indication of IP datagram fragmentation information.

IP Route Report

The IP Route Report contains an entry for each route presently known to this segment or node.

Dest.Address	Index	Metric1	Metric2	Metric3	Metric4	Metric5	Next Hop	Туре	Proto.	Age	Mask	Info
1.2.3.4	1	1	1	1	1	1	1.2.3.4	other	other	0:00:01.00	1.2.3.4	1.2.3
1.2.3.5	1	1	1	া	1	1	1.2.3.4	other	other	0:00:01.00	1.2.3.4	1,2,3
1.2.3.6	1	1	1	1	1	1	1.2.3.4	other	other	0:00:01.00	1.2.3.4	1.2.3
1.2.3.7	1	1	1	া	. 1	1	1.2.3.4	other	other	0:00:01.00	1.2.3.4	1.2.3
1.2.3.8	1	1	1	1	1	1	1.2.3.4	other	other	0:00:01.00	1.2.3.4	1.2.3
1.2.3.9	1	1	1	া	1	1	1.2.3.4	other	other	0:00:01.00	1.2.3.4	1,2,3
1.2.3.10	1	1	1	٦	1	1	1.2.3.4	other	other	0:00:01.00	1.2.3.4	1.2.3
1.2.3.11	1	1	1	্	1	1	1.2.3.4	other	other	0:00:01.00	1.2.3.4	1.2.3
1.2.3.12	1	1	1	1	1	1	1.2.3.4	other	other	0:00:01.00	1.2.3.4	1.2.3
1.2.3.13	1	1	1	1	1	1	1.2.3.4	other	other	0:00:01.00	1.2.3.4	1.2.3

CLEVER eRoute IP Route Reporting

There are five metrics provided by the IP Route report. The first metric is the primary routing parameter that identifies the routing-protocol specified in the router IP-Route-Proto value. If it is not used, its value should be set to -1. The second metric contains an alternate routing value that is determined by the IP-Route-Proto value. If it is not used, it should be set to -1. The third, fourth, and fifth metrics provide alternate routing information based on the value in IP-Route-Proto field.

The remaining fields are Next Hop, Type, Proto, Age, and Mask. The Next Hop field contains the IP address of the next hop of this route that the received packet can be sent to. The Type field provides the routing action. This field can have one of four possible values:

- 1 Indicates not applicable
- 2 Shows an invalidated route
- 3 Means to route only to the directly attached sub-network
- 4 Means to route to a non-local node/sub-network.

The Proto field is a routing mechanism that the route would be able to learn. There are four possible values:

- 1 Means no action
- 2 Means a non-protocol or manually configured entry
- 3 Indicates a network management protocol
- 4 Means to learn from ICMP data field.

The Age field contains the number of seconds since this route was last updated or otherwise determined to be correct. The Mask field provides the validation of the destination address via the logically ANDed comparison.

CONCLUDING REMARKS

TCP/IP route performance analysis can yield extremely beneficial data to better tune an open enterprise network environment, resulting in improved performance and reliability, as well as significant cost savings. It can be used with other TCP/IP performance analysis tools to effectively manage the performance of the TCP/IP network, particularly when it contains some TN3270-based networks. Given the nature of dynamic routing and its impact on problem diagnosis, performance tuning, and capacity planning, TCP/IP networks are understandably challenging to manage. Without the means to collect and analyze data in a systematic framework that shows the behavior of routes and segments over a period of time, symptoms may be easily misinterpreted. Thorough route performance analysis allows symptoms to be interpreted more accurately the first time, preventing the misuse of human resources by undertaking tasks that do not address the real problems. In addition, it can also prevent the unnecessary acquisition of additional hardware resources.

Through the description of data collection architectures, data elements, and examples, we hope we have demonstrated that superior data collection and analysis of TCP/IP route performance can indeed provide valuable input to address not only the short-term problems, but also to allow preparation to meet enterprise network long-term growth objectives. AES CLEVER *e*Route was created to uniquely answer to these needs.